

# IE605: Engineering Statistics

## Lecture 08

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## Previous Lecture:

- ▶ F-distributions
- ▶ Convergence of RVs
- ▶ Consistency
- ▶ Order Statistics

## This Lecture:

- ▶ Generating Random Samples

## Two methods

- ▶ Direct method
- ▶ Indirect method

### Direct method: Continuous case

- ▶ Generate random samples which has continuous CDF  $F$
- ▶  $U \sim Unif(0, 1)$ . Set  $X = F^{-1}(U)$ . Then  $X$  has CDF  $F$ .
- ▶  $P(X \leq x) = P(F^{-1}(U) \leq x) = P(U \leq F(x)) = F(x)$ .

**Example:** Generate samples  $\sim Exp(\lambda)$ .  $F(x) = 1 - e^{-\lambda x}$

- ▶  $F^{-1}(u) = -\frac{1}{\lambda} \log(1 - u)$
- ▶ Set  $X = -\frac{1}{\lambda} \log(1 - U) \sim Exp(\lambda)$

## Generating other continuous random variables

Gamma distribution  $Gamma(n, \lambda)$ : For some integer  $n$ ,  $\lambda > 0$

- ▶ For  $X_1, X_2, \dots, X_n$  are iid and  $\sim Exp(\lambda)$ , we know  $\sum_{i=1}^n X_i \sim Gamma(n, \lambda)$
- ▶  $U_1, U_2, \dots, U_n$  are iid. with  $\sim Unif(0, 1)$ .
- ▶  $X_i = -\frac{1}{\lambda} \log(1 - U_i)$ .  $X_i \sim Exp(\lambda)$
- ▶  $X = \sum_{i=1}^n X_i = -\frac{1}{\lambda} \sum_{i=1}^n \log(1 - U_i)$ .

Chi square distribution  $\chi_{2n}^2$ : For some integer  $n$

- ▶  $X \sim \chi_{2n}^2 = Gamma(n, 1/2)$

Beta distribution  $beta(m, n)$ : For some integers  $m, n$

- ▶  $U_1, U_2, \dots, U_m, U_{m+1}, \dots, U_{m+n}$  are iid  $\sim Unif(0, 1)$
- ▶  $X \sim \frac{\sum_{i=1}^m \log U_i}{\sum_{j=1}^{m+n} \log U_j} \sim beta(m, n)$ .

**Limitation:** We cannot generate Chi square method with odd degree of freedom. Particularly, Gaussian distribution!

## Other direct method

$$u = F(x) \implies u = \int_{-\infty}^x f(z) dz$$

To invert we need to solve integration for each  $u$ !

### Box-Muller method

Let  $U_1$  and  $U_2$  are iid with  $\sim Unif(0, 1)$ . Define

$$R = \sqrt{-2 \log U_1} \text{ and } \theta = 2\pi U_2.$$

Then

$$X_1 = R \cos(\theta) \text{ and } X_2 = R \sin(\theta)$$

are iid with  $\sim \mathcal{N}(0, 1)$ . (Verify!)

## Discrete case:

$X$  takes discrete values  $x_1, x_2, \dots$ . Assume  $x_1 < x_2 < x_3, \dots$ .  
To generate  $X$  satisfying CDF  $F$ .

$$P(F(x_i) < U \leq F(x_{i+1})) = F(x_{i+1}) - F(x_i) = P(x_{i+1}).$$

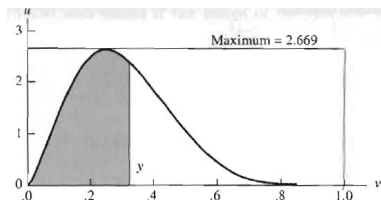
**Example:** Generate  $X \sim \text{Bin}(4, 5/8)$ .

Generate  $U \in \text{Unif}(0, 1)$  and set

$$X = \begin{cases} 0 & \text{if} & 0 < U \leq 0.02 \\ 1 & \text{if} & 0.02 < U \leq .152 \\ 2 & \text{if} & .152 < U \leq .481 \\ 3 & \text{if} & .481 < U \leq 0.847 \\ 4 & \text{if} & .847 < U \leq 1 \end{cases}$$

## Indirect method

Consider a distribution which takes value in  $[0, 1]$ .



Let  $(U, V)$  are iid  $\sim Unif(0, 1)$ .  $f$  is the given PDF.  $c = \max_y f(y)$

$$\begin{aligned} P\left(V \leq y, U \leq \frac{1}{c}f(V)\right) &= \int_0^y \int_0^{f_X(v)/c} dudv \\ &= \frac{1}{c} \int_0^y f_X(v)dv = \frac{1}{c}P(X \leq y) \end{aligned}$$

Set  $y = 1$ . We get  $P\left(U \leq \frac{1}{c}f(V)\right) = \frac{1}{c}$

## Indirect method contd..

$$\begin{aligned}P(X \leq y) &= \frac{P(V \leq y, U \leq \frac{1}{c}f(V))}{P(U \leq \frac{1}{c}f(V))} \\ &= P\left(V \leq y, U \leq \frac{1}{c}f(V) \mid U \leq \frac{1}{c}f(V)\right)\end{aligned}$$

### Algorithm:

1. (a) Generate  $(U, V)$  iid  $Unif(0, 1)$
2. if  $U \leq \frac{1}{c}f(V)$ . Set  $X = V$
3. else go to step 1.



## Accept/Reject Algorithm

Remove the requirement that distribution takes value in  $[0, 1]$ .

Let  $f$  be any given distribution and  $V$  is a RV with pdf  $f_V$

$$M = \max_y \frac{f(y)}{f_V(y)} < \infty (\text{assume!})$$

### Algorithm

1. Generate  $U \sim Unif(0, 1)$  and  $V \sim f_V$  independently
2. If  $U \leq \frac{1}{M}f(V)$ , set  $X = V$
3. Else, go to setp 1.

**Calim:**  $X$  has pdf  $f$ .

If  $M$  is not bounded . Accept/Reject algorithm will not work.  
Metropolis Algorithm extends the ideas.